

GLOBAL
EDITION



Concepts of Programming Languages

ELEVENTH EDITION

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ALWAYS LEARNING

PEARSON

Chapter 1 Preliminary

1.1 Reasons for Studying Concepts of Programming Languages

- Increased capacity to express ideas
- Improved background for choosing appropriate languages
- Increased ability to learn new languages
- Better understanding of the significance of implementation
- Better use of languages that are already known
- Overall advancement of computing

1.2 Programming Domains

- Scientific Applications
 - Fortran?
- Business Applications
 - COBOL (appeared in 1960)

1.2 Programming Domains

- Artificial Intelligence
 - **Symbolic** but not numeric
 - **Linked list** but not array
 - Functional language : LISP
 - Logic programming language: Prolog

1.2 Programming Domains

- Systems Programming
 - OS and tools
 - Machine-dependent
 - C language

1.2 Programming Domains

- Web Software
 - Markup languages
 - HTML, XML...
 - Scripting languages
 - JavaScript or PHP

1.3 Language Evaluation Criteria

- Impact on the software development process
- Maintenance

Readability

Maintenance and readability

- Overall Simplicity
 - Readability problems occur
 - Authors had learned a different subsets
 - Feature multiplicity
 - Operator overloading
 - Simplicity in languages can be carried too far
 - Result in less readable
 - Assembly language

Readability

- Orthogonality
 - Orthogonal
 - 直角的、正交的
 - Easier use (in mathematics)
 - Non-overlapping, uncorrelated, independent object
 - Definition of orthogonality in PL
 - First para. of Section 1.3.1.2

Readability

- Data Types
 - The presence of adequate facilities for defining data types and data structures in a language is another significant aid to readability

Readability

- Syntax Design
 - Special words
 - Form and meaning
 - Syntax and semantics

Writability

- Simplicity and Orthogonality
- Expressivity

Reliability

- Type checking
- Exception handling
 - Intercept run-time error
- Aliasing
 - A dangerous feature
 - E.g., Union & pointer in C
 - See next slice.
- Readability and Writability

Union of C

```
typedef struct a {  
    int i;  
    union {  
        float x;  
        int y; }  
} r1;  
r1.x=5.1;  
printf(“%d”,k.y);
```

Cost

- Cost of training programmers
- Cost of writing programs
- Cost of compiling programs
- Cost of executing programs
- Cost of the language implementation system
- Cost of poor reliability

Cost

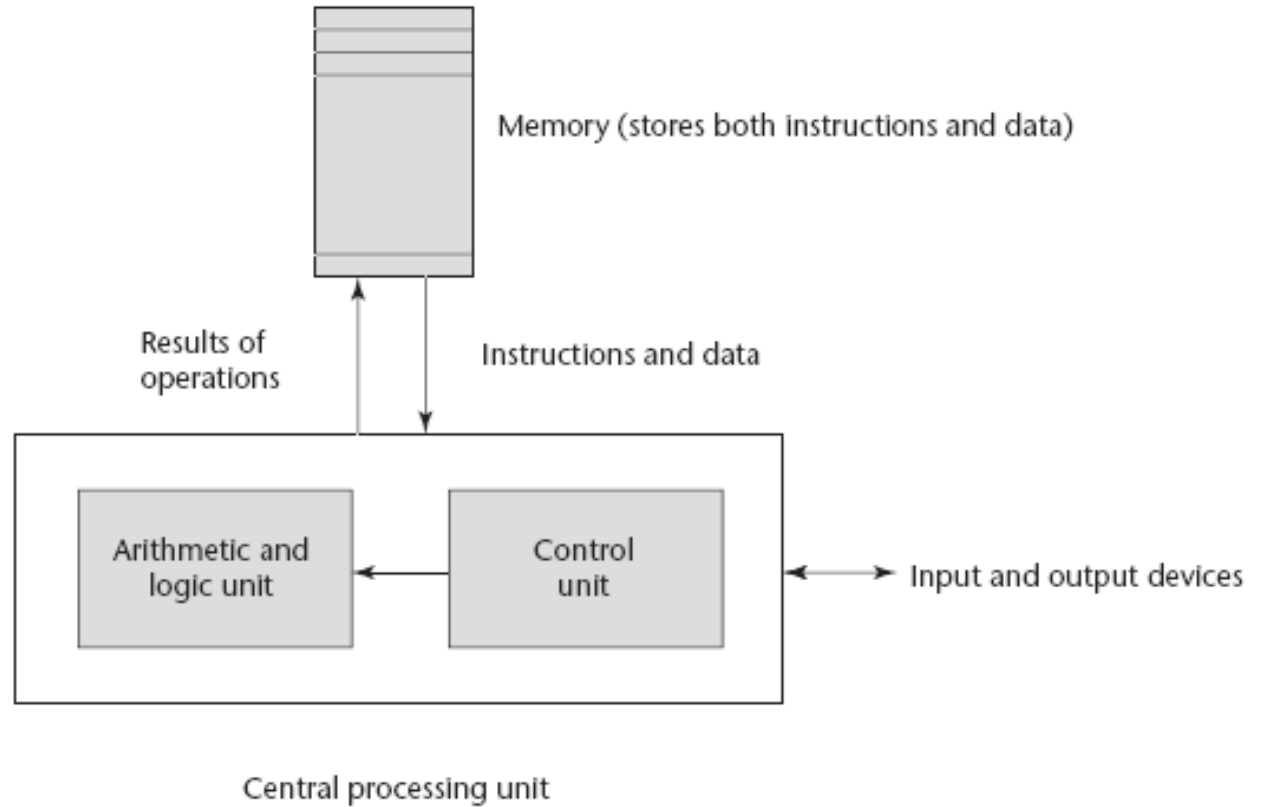
- Cost of maintaining programs
 - Maintenance costs can be as high as two to four times as much as development costs (Sommerville, 2005)
- Portability
- Generality and Well-definedness

1.4 Influences on Language Design

- Computer architecture
 - A profound effect on language design
 - Von Meumann architecture
 - Imperative languages
 - Central features
 - Variables
 - Assignment statements
 - Iterative form

Figure 1.1

The von Neumann computer architecture



1.4 Influences on Language Design

- Computer architecture
 - Languages that are not imperative
 - Functional language
 - Without assignment statements and without iteration
 - Imperative languages dominate!

1.4 Influences on Language Design

- Programming design methodologies
 - Trend
 - HW cost ↓
 - SW cost ↑

1.4 Influences on Language Design

- Programming design methodologies
 - 1970s
 - Top-down design and stepwise refinement
 - Late 1970s
 - A shift from procedure-oriented to data-oriented program design and methodologies
 - Abstract data types
 - 1980
 - Object-oriented design

1.5 Language Categories

- Four bins:
 - Imperative, functional, logic, and object-oriented.
- Others:
 - Scripting language
 - By interpretation
 - E.g., Perl, JavaScript, Ruby (still imperative)

1.5 Language Categories

- Recently days
 - Markup language
 - HTML, XML, XSLT, etc.

1.6 Language Design Trade-Offs

- What is the meaning of trade-off?
- Trade-offs
 - Reliability and cost of execution
 - Design trade-off
 - How about APL? (See next slice)
 - Writability and reliability

How does a really complicated APL routine look like?

Performing a fast Fourier transformation (FFT)

```
▽ Z←FFT X;C;D;E;J;K;LL;M;N;O
[1] LL←⌊2*-O-1M←⌊2⊕N,0ρE←1-2×~O←1J←1L
    ←0,0ρK←1N←1↑ρX
[2] →(M>L←L+1)/1+ρρJ←J,Nρ 0 1 0.=
    2*L)ρ 1
[3] Z←X[;(L←0)+(ϕLL)+.×J←(M,N)ρJ]
[4] X← 2 1 0.00(-O-K)÷1↑LL
[5] Z←Z[;K-,LL[L]×J[L;]]+(ρZ)ρ(-/X[;D]×
    Z[;C]),+/X[;D←O+NρLL[E+M-L]×-O-1
    2×LL[L]]×⊖Z[;C←K+,LL[L]×0=J[L;]]
[6] →((M+O)>L←L+1)/5
```

▽

1.7 Implementation Methods

- Compilations
- Pure Interpretation

Figure 1.4

Pure interpretation

